

Cauchy Fast Multipole Method for Analytic Kernel

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Outline

- 1 Overview
- 2 Algorithm
- 3 Results

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1 Overview

2 Algorithm

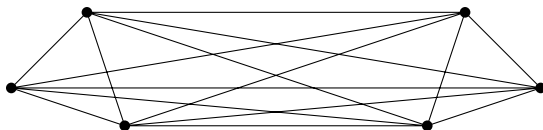
3 Results

Overview

Perform matrix-vector product of the form

$$\sum_{i=1}^N \sigma_i K(x_i, y_j) ; j = 1..N$$

e.g. electrostatic kernel : $K(x_i, y_j) = \frac{1}{|x_i - y_j|}$



Naive implementation : $O(N^2)$

Overview

One way to speed up the computation is to *group together* particles and use a *low-rank* approximation of the kernel.



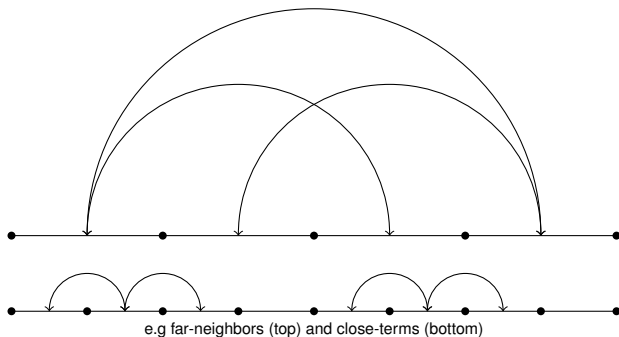
$$K(x, y) \approx \sum_{p=1}^P \sum_{q=1}^Q a_{pq} u_p(x) v_q(y)$$

$$\sum_{i=1}^N \sigma_i K(x_i, y_j) \approx \sum_{q=1}^Q \left(\sum_{p=1}^P a_{pq} \left(\sum_{i=1}^N \sigma_i u_p(x_i) \right) \right) v_q(y_j)$$

Tree decomposition

Usually, it is not possible to get an efficient *global* separable approximation. Instead, use approximations valid only in certain *region* (far apart) and use a *tree decomposition*.

- 1 Use different low-rank approximation at each scale/tree level
- 2 Compute close-term interactions at lowest level



Goal

We want to develop an algorithm with the following properties :

- 1 The algorithmic complexity is $O(N)$ (up to log-factor)
- 2 The translation and transfer operators are *diagonal*

$$K(x, y) \approx \sum_{p=1}^P \sum_{q=1}^Q a_{pq} u_p(x) v_q(y) \quad \text{vs} \quad \sum_{p=1}^P a_p u_p(x) v_p(y)$$

- 3 The algorithm is applicable to *any analytic kernel*

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Algorithm

The algorithm is based on *Cauchy Integral formula* i.e. given $K(z)$, $z \in \mathbb{C}$, complex differentiable $\forall z \in \Omega$ (holomorphic), then

$$K(x) = \frac{1}{2\pi i} \oint_{\Gamma} \frac{K(z)}{z - x} dz \quad (1)$$

Algorithm (Formal derivation)

Assume that $\Re(z - (x - y)) > 0 \forall z \in \Gamma$. Then, by the Laplace transform

$$K(x - y) = \frac{1}{2\pi i} \oint_{\Gamma} \frac{K(z)}{z - (x - y)} dz = \frac{1}{2\pi i} \oint_{\Gamma} K(z) \int_0^{\infty} e^{-s(z - (x - y))} ds dz$$

which upon discretization gives,

$$\begin{aligned} K(x - y) &\approx \frac{1}{2\pi i} \sum_n K(z_n) \sum_p e^{-s_p(z_n - (x - y))} \\ &= \sum_p \left(\frac{1}{2\pi i} \sum_n K(z_n) e^{-s_p z_n} \right) e^{s_p x} e^{-s_p y} \end{aligned}$$

This expansion possesses the desired form

Algorithm (Formal derivation)

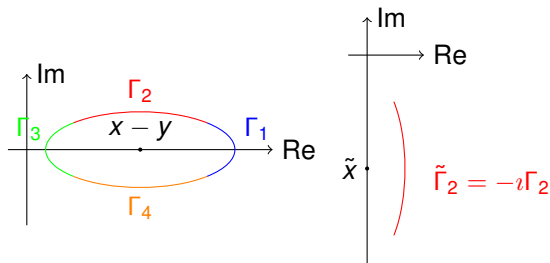
Current issues with the separable expansion :

- 1 It is not true that $\Re(z - (x - y)) > 0 \forall z \in \Gamma$.
- 2 The use of exponentials $e^{-s_k z_n}$, $e^{s_k x}$ and $e^{-s_k y}$ is numerically unstable

Both problems can however be solved.

Algorithm (Formal derivation)

To solve the first problem, we break the contour into paths and rotate each path by multiplying by $\{\lambda_k\} \in \{1, -i, -1, i\}$.



$$\frac{1}{2\pi i} \int_{\Gamma_k} \frac{\lambda_k}{\lambda_k} \frac{K(z)}{z - (x-y)} dz = \frac{\lambda_k}{2\pi i} \int_{\Gamma_k} K(z) \int_0^\infty e^{-s\lambda_k(z-(x-y))} ds dz$$

Algorithm (Formal derivation)

To solve the stability problem, we introduce parameters $\{\ell_k\}$ such that

$$\Re(\lambda_k z - 2\ell_k) > 0 \quad (2)$$

$$\Re(\ell_k - \lambda_k x) > 0 \quad (3)$$

$$\Re(\ell_k + \lambda_k y) > 0 \quad (4)$$

Thanks to the well-separated condition, the tree decomposition and the choice of paths, it is **always possible** to find such parameters.

Then,

$$\int_0^\infty e^{-s\lambda_k(z-(x-y))} ds = \int_0^\infty e^{-s(\lambda_k z - 2\ell_k)} e^{-s(\ell_k - \lambda_k x)} e^{-s(\ell_k + \lambda_k y)} ds$$

which can be discretized efficiently to produce a **stable** approximation.

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Results

Kernel studied,

- Algebraic growth (Multiquadric) : $K(x - y) = \sqrt{1 + (x - y)^2}$
- Algebraic decay (Inverse Multiquadric) : $K(x - y) = \frac{1}{\sqrt{1+(x-y)^2}}$
- Oscillatory (Helmholtz) : $K(x - y) = \frac{e^{ik|x-y|}}{k|x-y|}$

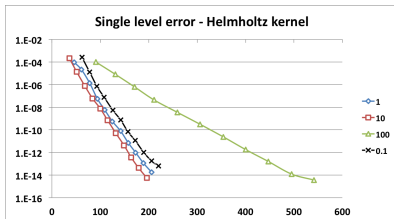
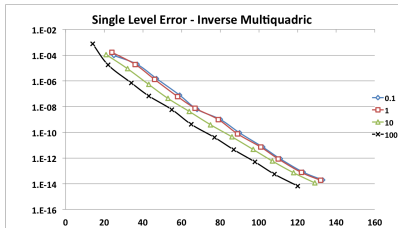
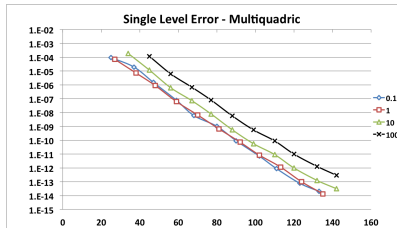
Error considered,

- Maximum absolute error : $\max_n |\phi_n^{FMM} - \phi_n^{Direct}|$
- Relative L^2 error : $\frac{\sqrt{\sum_n (\phi_n^{FMM} - \phi_n^{Direct})^2}}{\sqrt{\sum_n (\phi_n^{Direct})^2}}$

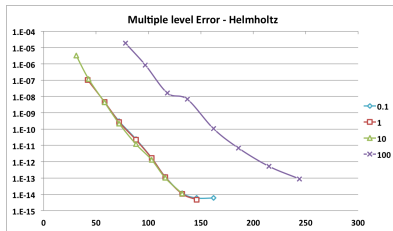
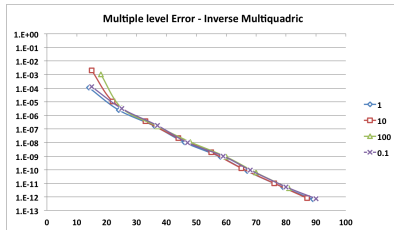
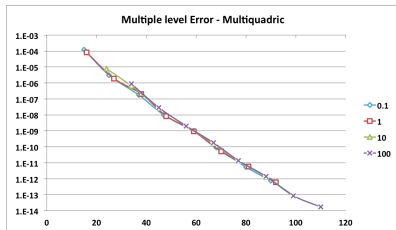
Particle distribution,

- N locations *uniformly* distributed in an interval
- Various interval size : 0.1, 1, 10, 100

Results (Single level - Maximum absolute error vs # Quadrature points)

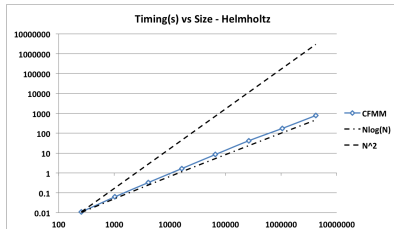
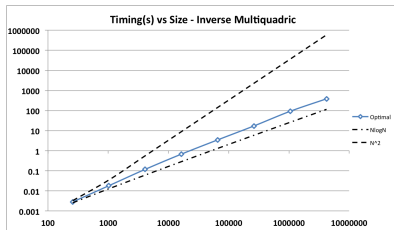
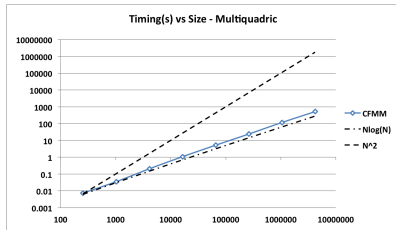


Results (Multiple levels - Relative L^2 error vs # Quadrature points)



Number of particles : 65536
Number of levels : 8

Results Multiple levels - Timing vs Problem Size



Interval length : 100

Relative L^2 error : Multiquadric(10^{-6}), Inverse Multiquadric(10^{-3}), Helmholtz(10^{-5})

Computer : Dell 3GHz CPU, 8GBytes RAM

Conclusion

- 1 The algorithmic complexity is $O(N \log(N))$
- 2 The translation and transfer operators are *diagonal*
- 3 The algorithm is applicable to *any analytic kernel*
- 4 Extension to n-D is trivial

$$K(x, y, z) = \frac{1}{(2\pi i)^3} \oint_{\Gamma_1} \oint_{\Gamma_2} \oint_{\Gamma_3} \frac{K(z_1, z_2, z_3)}{(z_1 - x)(z_2 - y)(z_3 - z)} dz_1 dz_2 dz_3$$